Hecke group algebras as degenerate affine Hecke algebras

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It all started at FPSAC 2006, San Diego



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Coxeter groups

Definition (Coxeter group W)

Generators : $(s_i)_{i \in S}$ (simple reflections)

Relations: $s_i^2 = 1$ and $\underbrace{s_i s_j \cdots}_{m_{i,i}} = \underbrace{s_j s_i \cdots}_{m_{i,i}}$, for $i \neq j$

Group algebra: $\mathbb{C}[W]$

Example (Type A_n : symmetric group \mathfrak{S}_{n+1})

Generators: $(s_i)_{i=1,...,n}$ (elementary transpositions) Relations:

$$egin{aligned} s_i^2 &= 1 & & ext{for all } 1 \leq i \leq n, \ s_i s_j &= s_j s_i & ext{for all } |i-j| > 1, \ s_{i+1} s_i &= s_{i+1} s_i s_{i+1} & ext{for all } 1 \leq i \leq n-1 \end{aligned}$$

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0-(Iwahori)-Hecke algebras

Definition (0-Hecke algebra H(W)(0))

Generators : $(\pi_i)_{i \in S}$

Relations: $\pi_i^2 = \pi_i$ and $\underbrace{\pi_i \pi_j \cdots}_{m_{i,i}} = \underbrace{\pi_j \pi_i \cdots}_{m_{i,i}}$ for $i \neq j$

Basis: $(\pi_w)_{w \in W}$

Example (Type A_n)

Generators: $(\pi_i)_{i=1,...,n}$ (adjacent comparators)

Relations:

$$\pi_i^2 = \pi_i$$
 for all $1 \le i \le n$, $\pi_i \pi_j = \pi_j \pi_i$ for all $|i - j| > 1$, $\pi_{i+1} \pi_i = \pi_{i+1} \pi_i \pi_{i+1}$ for all $1 \le i \le n-1$

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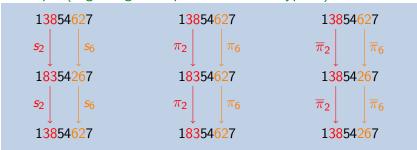
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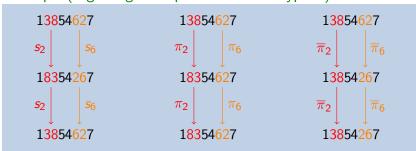
$$\begin{split} \pi_i^2 &= \pi_i & \text{for all } 1 \leq i \leq n, \\ \pi_i \pi_j &= \pi_j \pi_i & \text{for all } |i-j| > 1, \\ \pi_i \pi_{i+1} \pi_i &= \pi_{i+1} \pi_i \pi_{i+1} & \text{for all } 1 \leq i \leq n-1. \end{split}$$

Example (Right regular representation for type A)



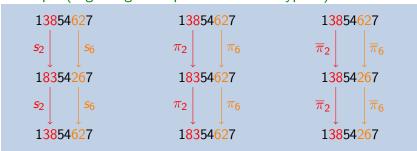
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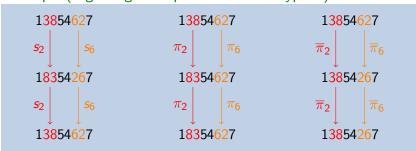
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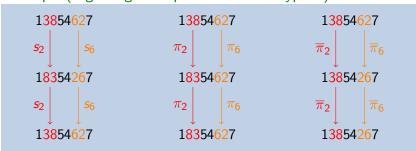
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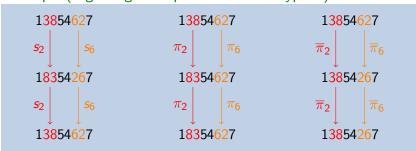
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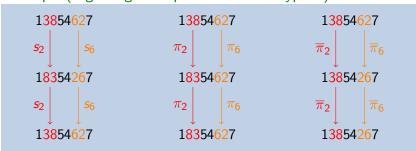
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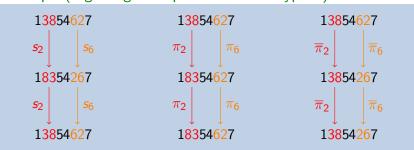
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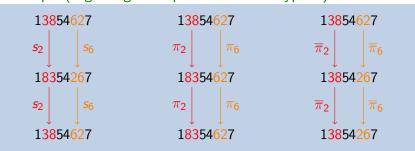
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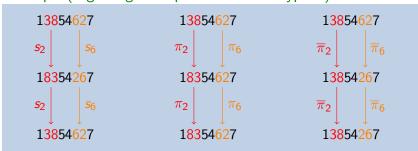
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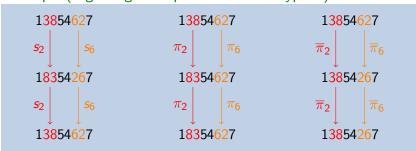
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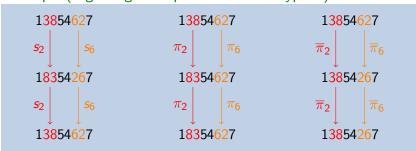
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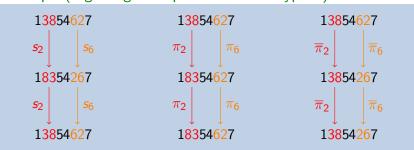
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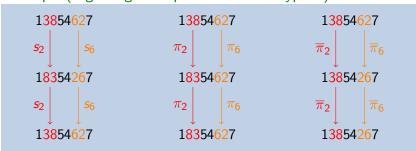
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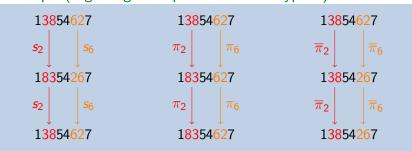
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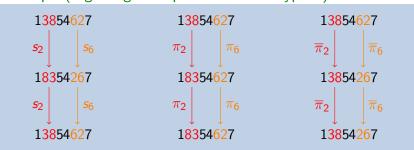
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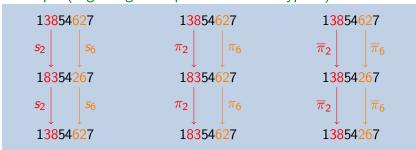
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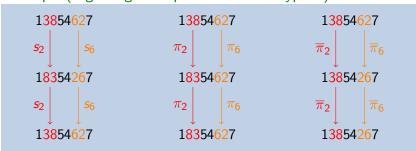
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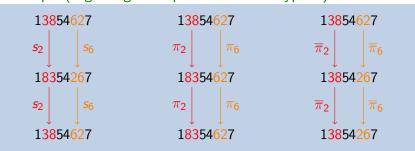
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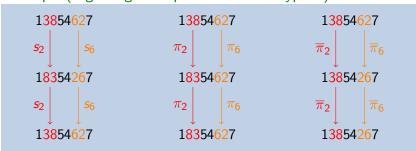
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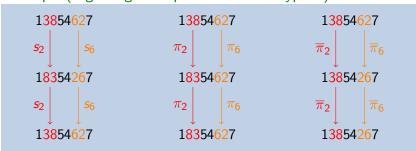
876<mark>5</mark>1342

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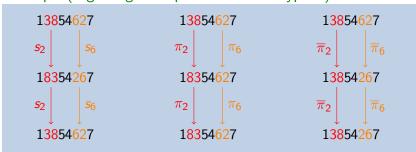
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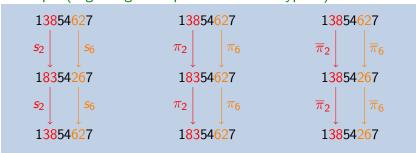
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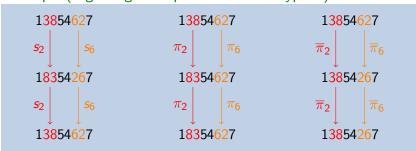
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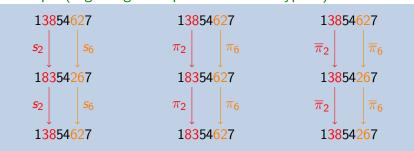
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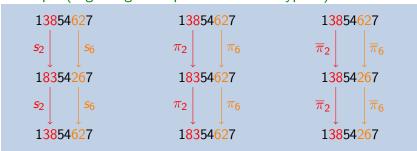
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Hecke algebras

Take q_1 and q_2 parameters, and set $q:=-\frac{q_1}{q_2}$.

Definition (Hecke algebra $H(W)(q_1, q_2)$)

Generators : $(T_i)_{i \in S}$ Relations: $(T_i - q_1)(T_i - q_2) = 0$ and $\underbrace{T_i T_j \cdots}_{m_{i,j}} = \underbrace{T_j T_i \cdots}_{m_{i,j}}$, for $i \neq j$ Basis: $(T_w)_{w \in W}$

- At q=1: group algebra $\mathbb{C}[W]$
- At q = 0: 0-Hecke algebra H(W)(0)
- At q not 0 nor a root of unity: isomorphic to $\mathbb{C}[W]$

Realization of T_i as operator in End($\mathbb{C}W$):

$$T_i := (q_1 + q_2)\pi_i - q_1s$$

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A silly idea during a brainstorm (Thibon, Novelli, H., T., 2003)

Definition (Hecke group algebra HW of a Coxeter group W)

$$HW := \langle (\pi_i, s_i)_{i \in S} \rangle \subset \operatorname{End}(\mathbb{C}W)$$

- Any interesting structure?
- Contains all Hecke algebras by construction
- Type A: dimension and dimension of the radical in the Sloane!

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The Hecke group algebra of rank 1

$$\begin{aligned} \mathcal{W} &:= \{1,s\} & \quad \mathbb{C}\mathcal{W} := \mathbb{C}.1 \oplus \mathbb{C}.s \\ \mathrm{id} &= \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, \quad s = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \quad \pi = \begin{pmatrix} 0 & 0 \\ 1 & 1 \end{pmatrix}, \quad \overline{\pi} = \begin{pmatrix} 1 & 1 \\ 0 & 0 \end{pmatrix} \end{aligned}$$

$$s\pi=\pi, \qquad \pi s=\overline{\pi}, \qquad \overline{\pi}+\pi=1+\overline{\pi}$$
 $\{\mathrm{id},s,\pi\} \qquad \mathrm{or} \qquad \{\mathrm{id},\pi,\overline{\pi}\}$

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Theorem (H., T., 2005)

- HW algebra of left antisymmetry preserving operators
- HW* algebra of left symmetry preserving operators
- Basis of HW: $\{w\pi_{w'} \mid \mathsf{D}_R(w) \cap \mathsf{D}_L(w') = \emptyset\}$
- Rep. theory governed by the combinatorics of descents
- HW Morita equivalent to the poset algebra of boolean lattice
- Projective & simple modules indexed by parabolic subgroups Restriction of simple:
 - Exactly the Young's ribbon representation of W
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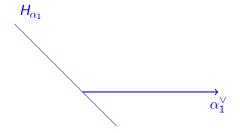
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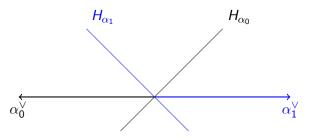
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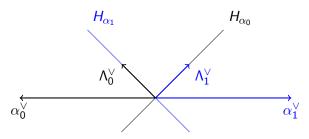
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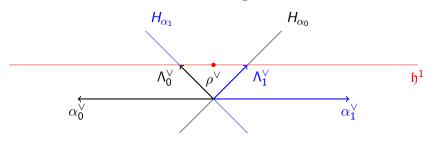


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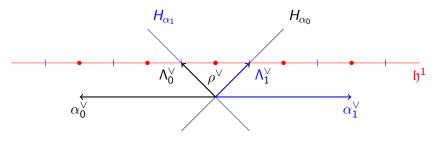


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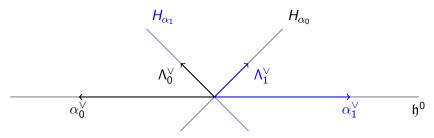
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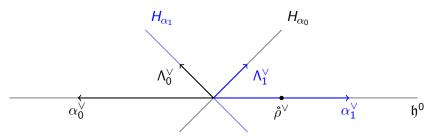
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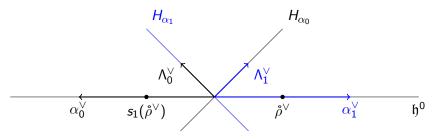
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$$u = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, \quad \pi_1 = \begin{pmatrix} 0 & 0 \\ 1 & 1 \end{pmatrix}, \quad \pi_0 = \begin{pmatrix} 1 & 1 \\ 0 & 0 \end{pmatrix}$$

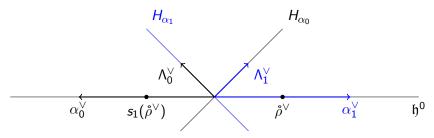
- π_0, π_1 acts transitively on \mathring{W}
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Remark (At level 0)

$$m = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, \quad \pi_1 = \begin{pmatrix} 0 & 0 \\ 1 & 1 \end{pmatrix}, \quad \pi_0 = \begin{pmatrix} 1 & 1 \\ 0 & 0 \end{pmatrix}$$

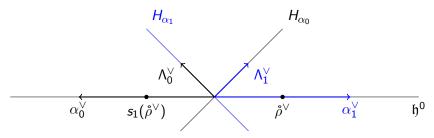
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Remark (At level 0)

$$\mathrm{id} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, \quad \pi_1 = \begin{pmatrix} 0 & 0 \\ 1 & 1 \end{pmatrix}, \quad \pi_0 = \begin{pmatrix} 1 & 1 \\ 0 & 0 \end{pmatrix}$$

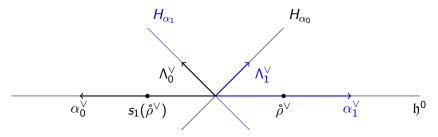
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Theorem (H.,S.,T. 2008)

W: affine Weyl group

W: finite Weyl group induced by the level 0 action

- W degenerates trivially to W
- H(W)(0) degenerates to $H\mathring{W}$
- H(W)(q) degenerates to $H\mathring{W}$, for q generic
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 π_1, \ldots, π_n can antisort 12345 to 54321 (bubble sort) But not back! (going down the permutohedron)

Definition (Affine action)

Put the permutation on a circle π_0 acts between last and first letter

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Proposition (S., T., 2007)

- Similar algorithms for types B, C, D
- Existence for all types (including twisted)

• Type B:
$$0 \ge 2 - 3 \implies 4$$
 $1 < 2 < 3 < 4 < \underline{4} < \underline{3} < \underline{2} < \underline{1}$

- Type-free induction strategy
- Case by case induction step

 Brute force on computer for the exceptional types (*E*₈!)

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Proof

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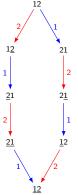
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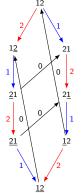
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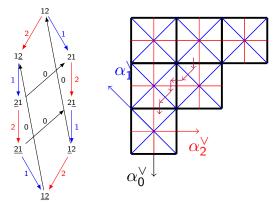
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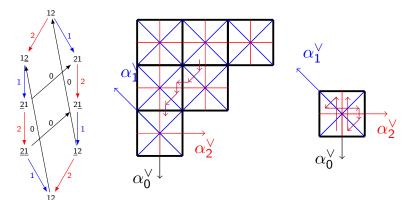
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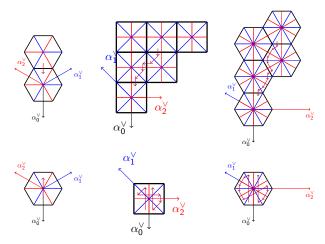
 π_0, π_1, π_2 on C_2 Alcove picture at level 1



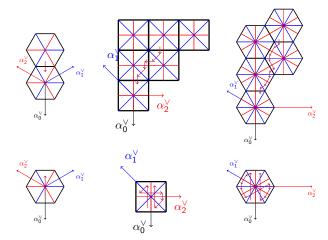
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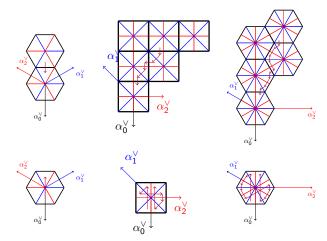
Quotient at level 0 (Steinberg torus)



- Covers all rank 2 affine Weyl groups
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Theorem (S.,T., 2008)

```
Take q non zero and non root of unity, and t:Y^{\lambda^\vee}\mapsto q^{-\operatorname{ht}(\lambda^\vee)}
Then, 
ho_t(\mathsf{H}(W)(q))=\mathsf{H}\mathring{W}
I.e. \mathsf{H}\mathring{W} (non trivial!) quotient of \mathcal{H}(q,t) and of \mathsf{H}(W)(q)
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Conclusion

- Hecke group algebras:
 - Many equivalent definitions
 - Nice structure and representation theory
 - (type A) Connections with NCSF, parking functions
 - Connections with 0-Hecke and affine Hecke algebras
- Where does this structure come from?
- Is it useful?